

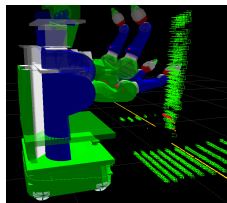
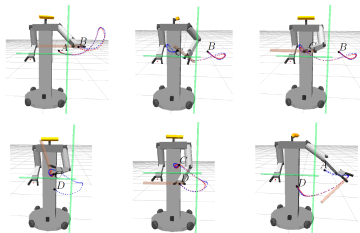
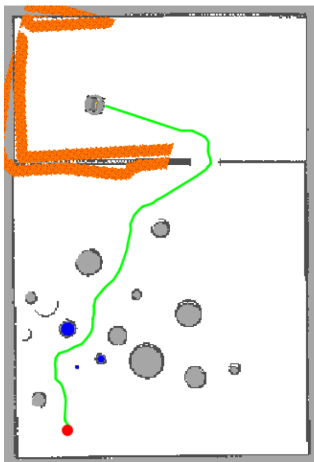
Motion Planning and Control for Domestic Service Robots

J.J.M. Lunenburg

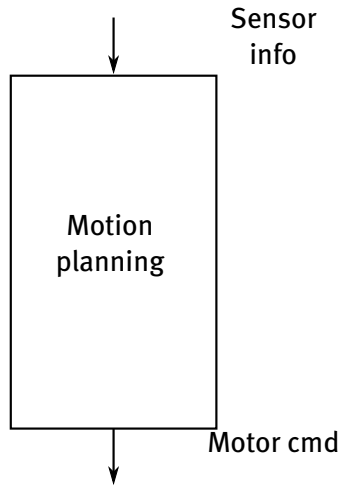
May 8, 2014

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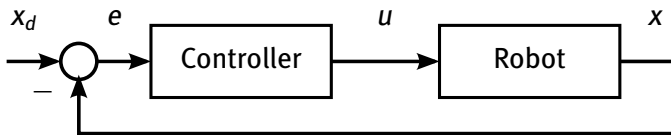
Where innovation starts



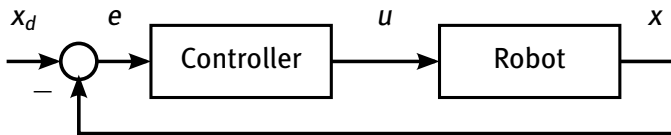
- ▶ **Sensors:**
 - Laser scan
 - Odometry
- ▶ **Motor cmd**
 - Velocity in x , y and ϕ direction
- ▶ Move towards a desired pose
- ▶ Don't crash into the wall!



- ▶ Feedback control!
- ▶ Carrot planner
 - P(D)-controller
- ▶ Dynamic Window Approach:
 - Search for a translational and rotational velocity
 - Optimization over a finite horizon
 - MPC-controller



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- ▶ **What about obstacles?**

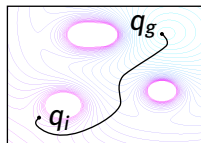
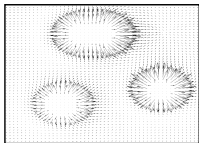
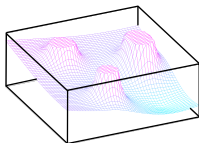
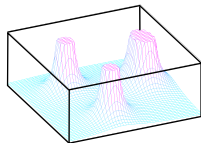
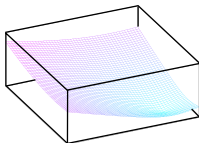
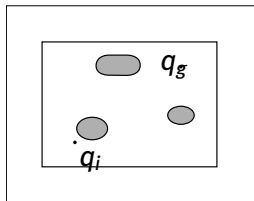


Don't crash into obstacles

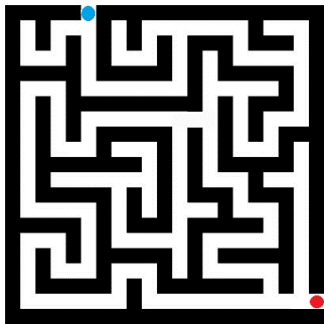
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- ▶ Line collision checks
- ▶ Forward simulation: rejects inputs

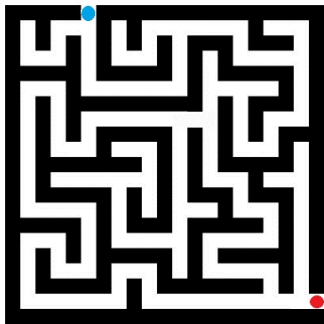
- ▶ Line collision checks
- ▶ Forward simulation: rejects inputs
- ▶ Alternative: potential fields
 - Goal and obstacles form attractive and repulsive forces



- ▶ Local methods
 - Vicinity of the robot
- ▶ Completeness: ‘getting stuck’
- ▶ Optimality: ‘the shortest path’



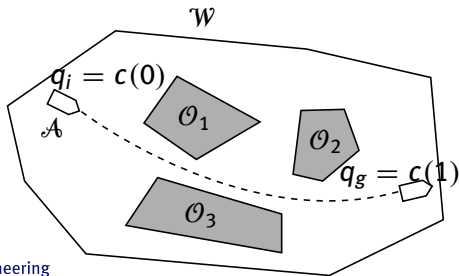
- ▶ Local methods
 - Vicinity of the robot
- ▶ Completeness: 'getting stuck'
- ▶ Optimality: 'the shortest path'
- ▶ Global connectivity information required



With:

- ▶ Pose (position and orientation)
- ▶ Single rigid body \mathcal{A}
- ▶ n-dimensional Euclidean space $\mathcal{W} = \mathbb{R}^n$
- ▶ Static, rigid obstacles \mathcal{O}_i in \mathcal{W}

Given an initial pose and a goal pose of \mathcal{A} in \mathcal{W} , find a path c in the form of a continuous sequence of poses of \mathcal{A} that do not collide or contact with \mathcal{O}_i , that will allow \mathcal{A} to move from its starting pose to its goal pose and report failure if such a path does not exist.



Six specifications and properties

- ▶ **Completeness:** finding a path if one exists
- ▶ **Optimality:** finding the optimal path
- ▶ **Computational complexity**
- ▶ **Robustness against a dynamic environment**
- ▶ **Robustness against uncertainty**
- ▶ **Kinematic and dynamic constraints**

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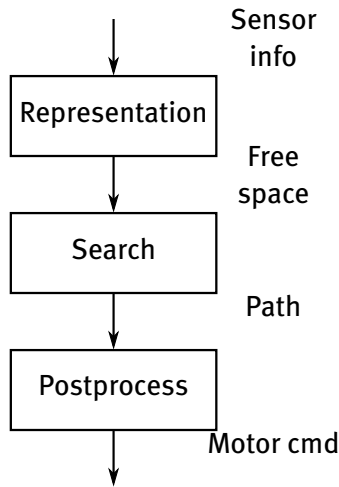
So how do we approach this problem?

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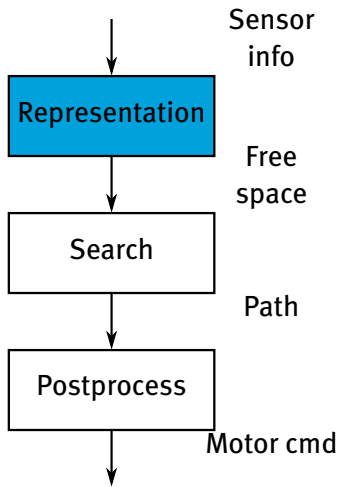
So how do we approach this problem?

Representation and searching!

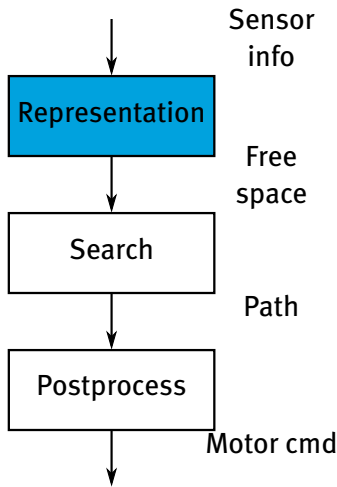


► The configuration space

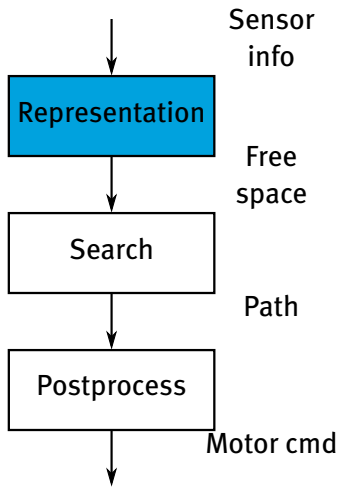
- Simplifies the problem: search for a solution for a single point
- Generic
- Computationally efficient

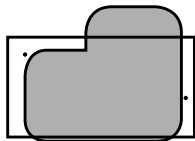
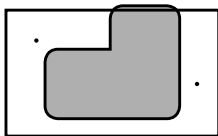
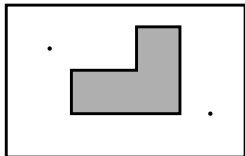
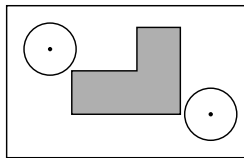
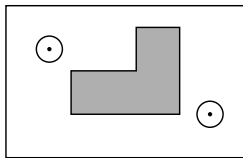
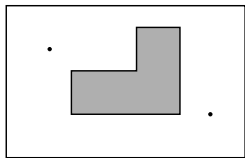


- ▶ The configuration space
 - Simplifies the problem: search for a solution for a single point
 - Generic
 - Computationally efficient
- ▶ Representation methods:
 - Exact
 - Roadmaps
 - Exact cell decomposition
 - Approximate
 - Approximate cell decompositions
 - Sampling-based methods
 - Potential fields



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- ▶ Common assumption: localization

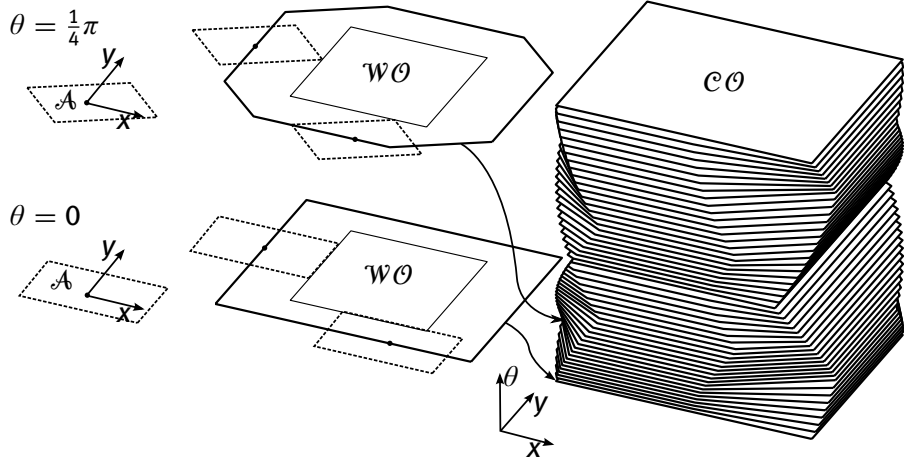




(a)

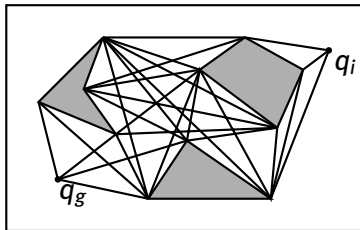
(b)

(c)



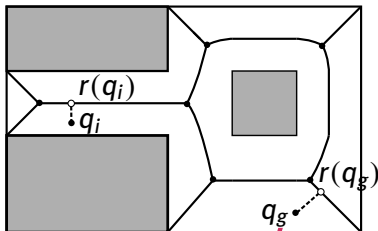
Visibility graph

- ▶ Two nodes are connected if the straight line between them is collision-free
- ▶ $\dim(\mathcal{C}) \leq 2$
- ▶ Optimal w.r.t. distance traveled



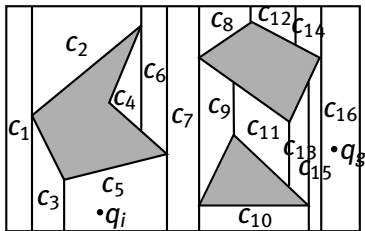
Deformation retracts

- ▶ ‘Shrink’ a space into a subspace
- ▶ (Generalized) Voronoi diagram
- ▶ Optimal w.r.t. distance to obstacles



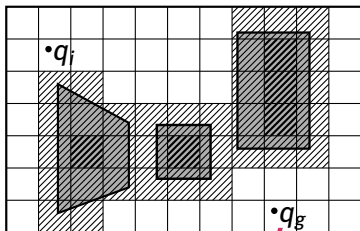
Exact decomposition

- ▶ Trapezoidal decomposition
- ▶ Sweep line algorithm
- ▶ Non-optimal



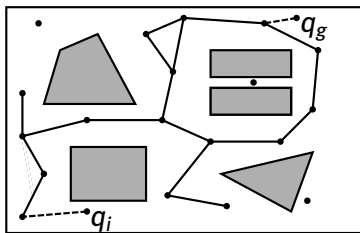
Approximate decomposition

- ▶ Obstacle boundaries do not coincide with cell boundaries
- ▶ Free cells, mixed cells and occupied cells
- ▶ Resolution complete



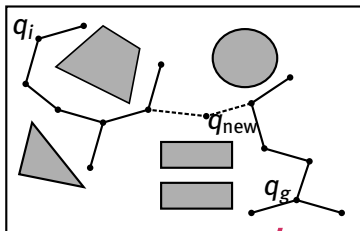
Probabilistic roadmap

- ▶ Learning phase: sample configuration q_{rand} and check for collisions
- ▶ Query phase: connect q_i and q_g to roadmap \mathcal{R}
- ▶ Probabilistically complete

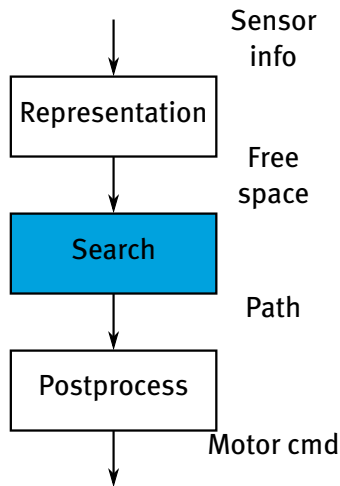


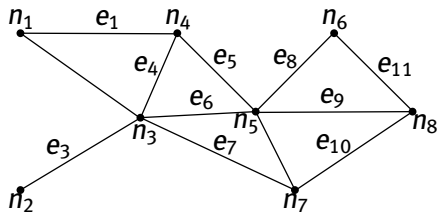
Single-query planner

- ▶ Explore relevant subset of \mathcal{C}_{free}
- ▶ (Bidirectional) Rapidly-exploring Random Tree
- ▶ No search algorithm required
- ▶ Probabilistically complete, non-optimal

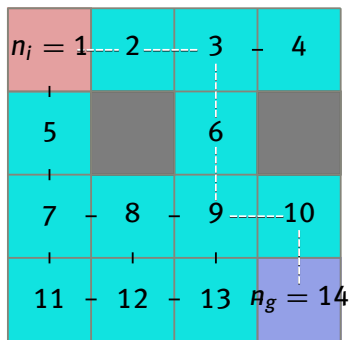


- ▶ Graphs and costmaps
- ▶ Graph search algorithms:
 - Uninformed
 - Informed
 - Local

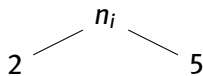
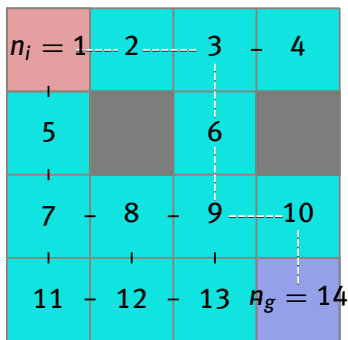




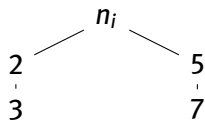
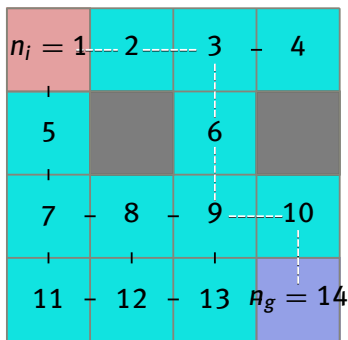
- ▶ Nodes (vertices) and edges
- ▶ Including weights: costmap
- ▶ Parent: node with subsequent nodes (children)
- ▶ Branch: series of nodes connecting the root to a leaf
- ▶ Frontier: set of all leaf nodes available for expansion
- ▶ Closed list (explored set): nodes that have been visited
- ▶ Expansion is determined by function $f(n)$

 n_i

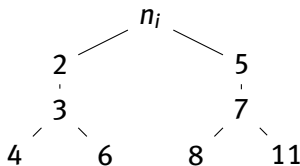
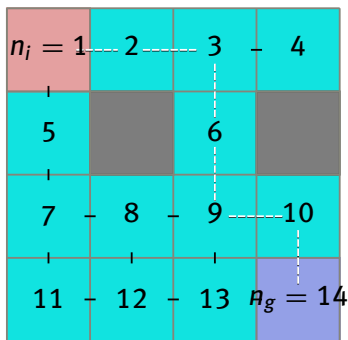
- ▶ $f(n) = g(n)$, with $g(n)$ a FIFO queue
- ▶ All nodes at a certain depth are expanded before going to the next level
- ▶ Complete (if 'branching' factor is finite)
- ▶ **Optimal**: only if all edges have equal costs



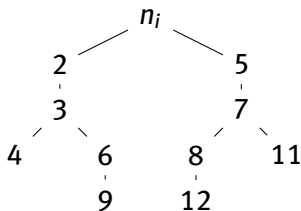
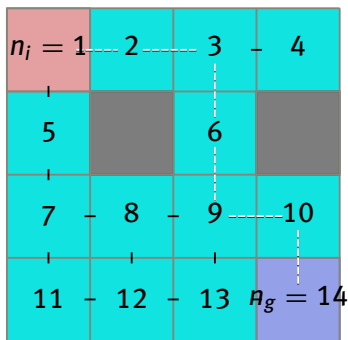
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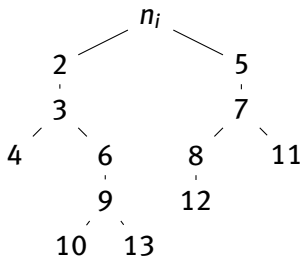
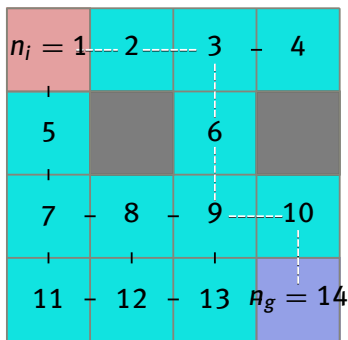
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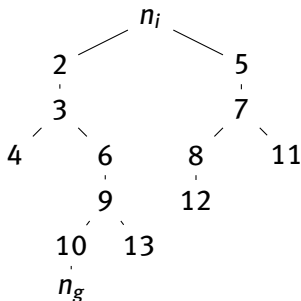
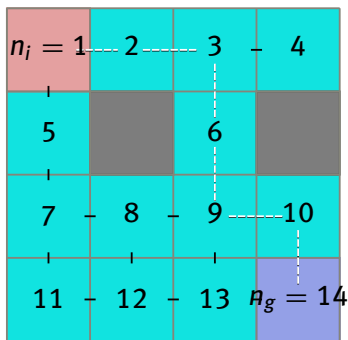
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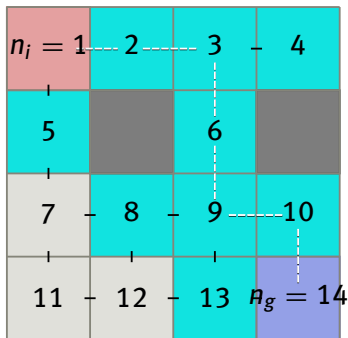
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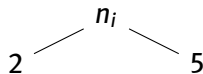
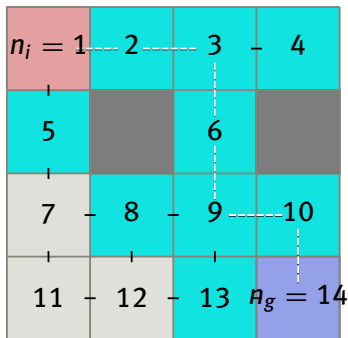


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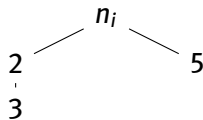
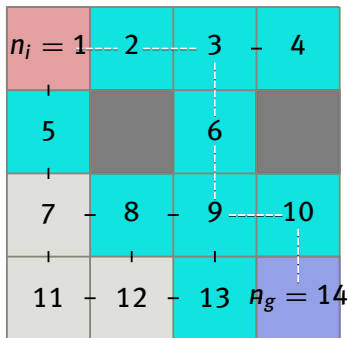
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- ▶ The most recent expanded node is put the beginning of the stack
- ▶ Completeness: if search space is finite
- ▶ Not optimal

- Example: put goal at node 5



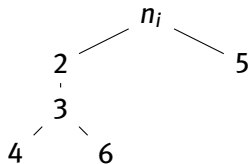
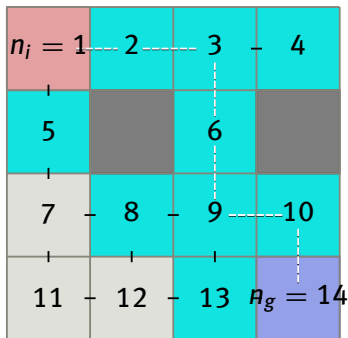
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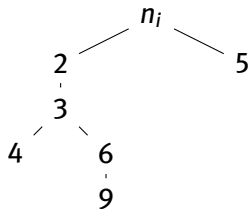
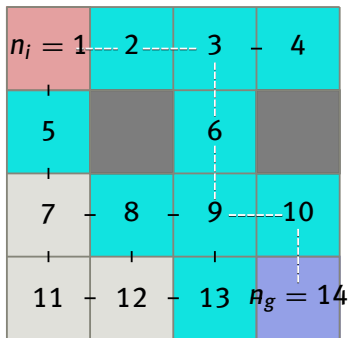
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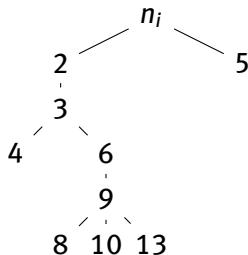
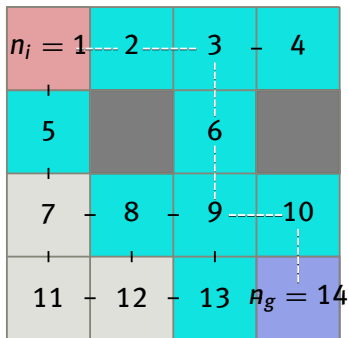
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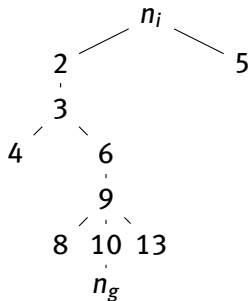
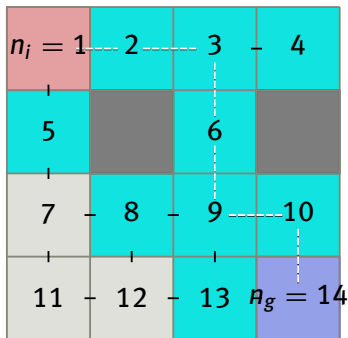
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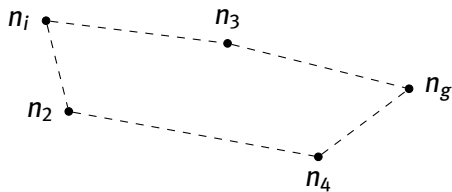
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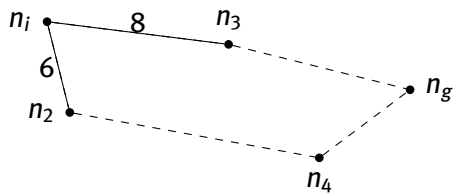
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- Example: put goal at node 5

- ▶ $f(n) = g(n)$, with $g(n)$ a priority queue
- ▶ The node with the lowest cost is expanded
- ▶ Completeness: if search space is finite
- ▶ Optimal

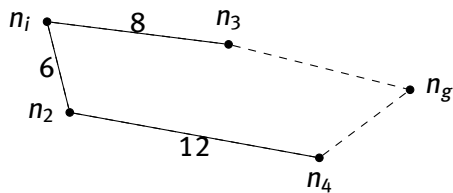


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- ▶ Optimal



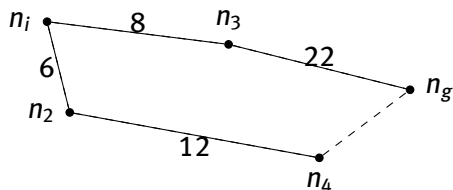
- ▶ $6 < 8 \rightarrow$ expand n_2

- ▶ $f(n) = g(n)$, with $g(n)$ a priority queue
- ▶ The node with the lowest cost is expanded
- ▶ Completeness: if search space is finite
- ▶ Optimal



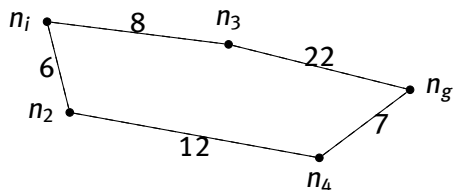
- ▶ $6 < 8 \rightarrow$ expand n_2
- ▶ $8 < 6 + 12 \rightarrow$ expand n_3

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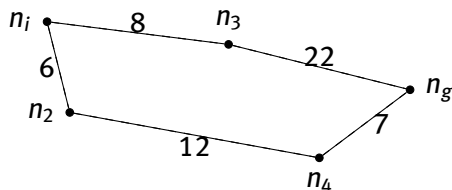
- ▶ $6 < 8 \rightarrow$ expand n_2
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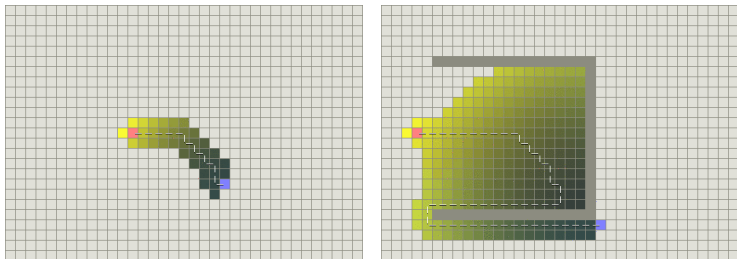
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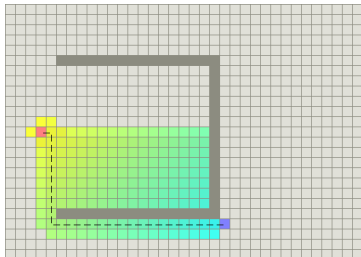
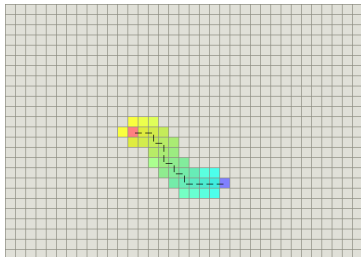


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- ▶ $8 < 6 + 12 \rightarrow$ expand n_3
- ▶ n_g reached, but $8 + 22 > 6 + 12$

Why not use knowledge of the goal location?

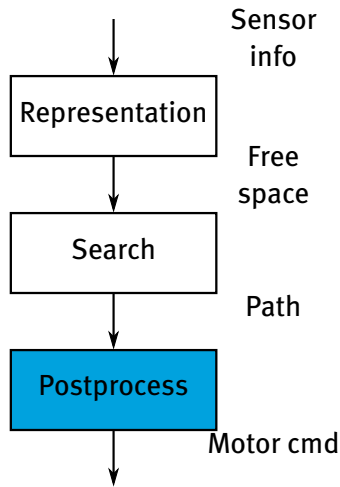


- ▶ $f(n) = h(n)$, with $h(n)$ a heuristic (distance) function
- ▶ Expands the node closest to the goal
- ▶ Complete
- ▶ Non-optimal (see figure)

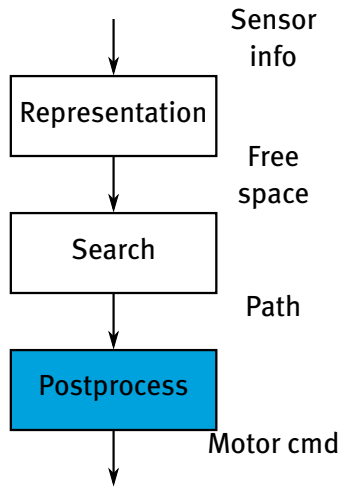
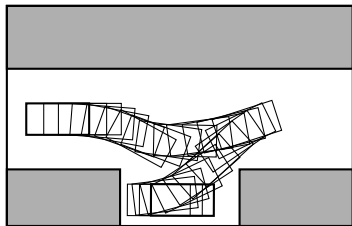


- ▶ $f(n) = g(n) + h(n)$, with $g(n)$ costs to reach a node and $h(n)$ heuristic to reach the goal
- ▶ Takes both costs into account
- ▶ Complete
- ▶ Optimal if the heuristic function is consistent:
 - $h(n) \leq c(n \rightarrow n') + h(n')$

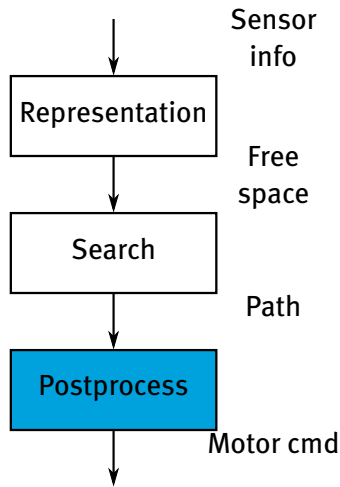
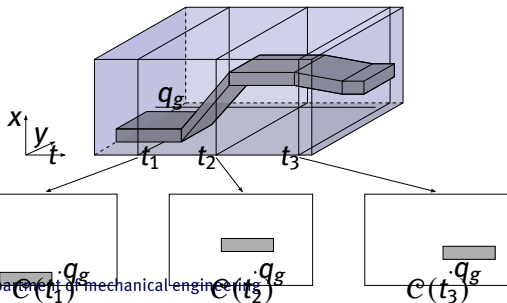
- ▶ The path resulting from searching the representation is not yet suitable for execution



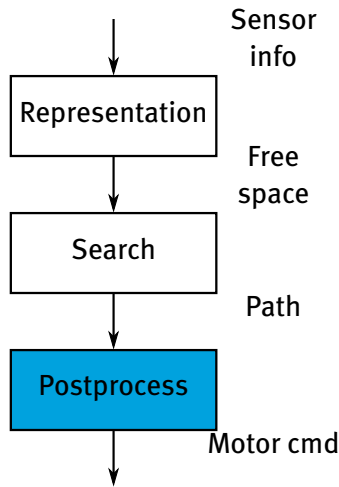
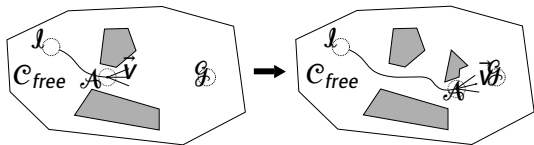
- ▶ The path resulting from searching the representation is not yet suitable for execution
- ▶ Kinodynamic constraints



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- ▶ Kinodynamic constraints
- ▶ Dynamic environments



- ▶ The path resulting from searching the representation is not yet suitable for execution
- ▶ Kinodynamic constraints
- ▶ Dynamic environments
- ▶ Uncertainty



Decoupled trajectory planning

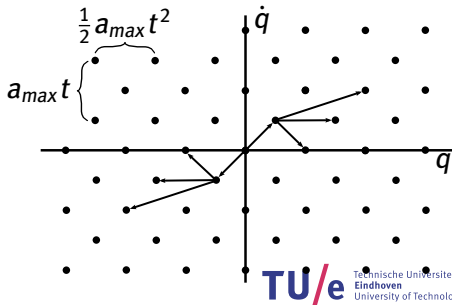
- ▶ Path planning: collision free path c in $\mathcal{C}_{\text{free}}$
- ▶ Transform c into c' , satisfying non-holonomic constraints
- ▶ Compute timing function such that $c'(t)$ satisfies kinodynamic constraints

Decoupled trajectory planning

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Direct trajectory planning

- ▶ Searching on a lattice
- ▶ Sampling based methods: select input at random from set of admissible controls



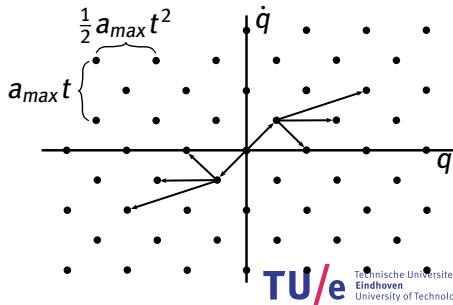
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Motion primitives



Re-planning (of an entire path)

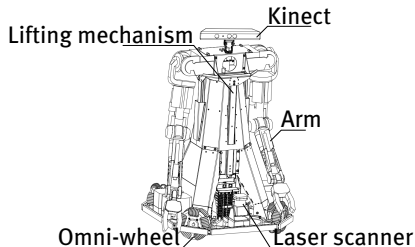
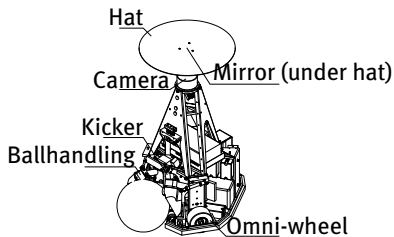
- ▶ Re-planning from the current situation
- ▶ Reuse information of previous searches (incremental search)
- ▶ The planner can return an (approximate and suboptimal) plan at any time (anytime planning)

- ▶ Reduction of complexity: divide the planning problem into global and local planner
 - Global planner: computes a path from start to goal
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- ▶ Topological maps
 - Abstract representation that describes relationships between features of the environment
 - Compact and stable w.r.t. sensor noise and uncertainty

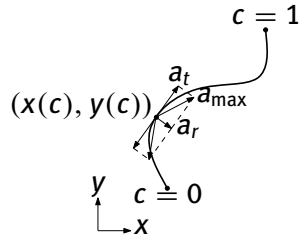
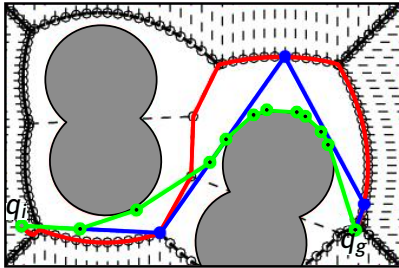
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How is motion planning applied in TU/e?



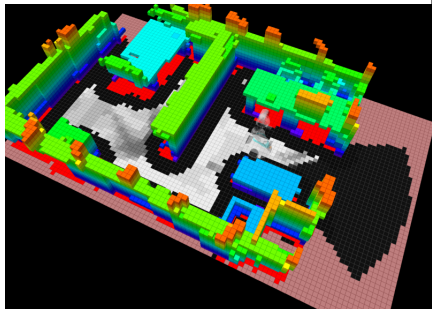
- ▶ Soccer pitch
- ▶ 12 m × 18 m
- ▶ Known environment
- ▶ Dynamic obstacles (hostile)
- ▶ 3 m/s

- ▶ House/care environment
- ▶ Arbitrary size
- ▶ Partially unknown
- ▶ Static and dynamic obstacles
- ▶ 1 m/s

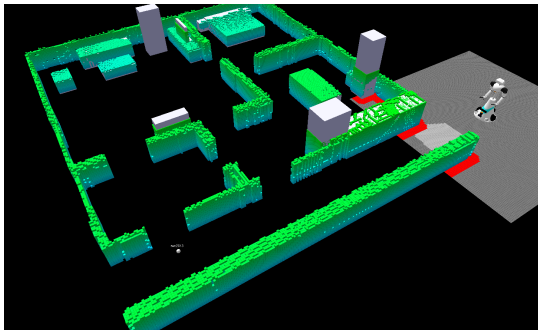


- ▶ Voronoi diagram representation, searched with Dijkstra's algorithm
- ▶ Shortcut algorithm to cut-off sharp turns
- ▶ Time-optimal trajectory through waypoints using Bézier curves

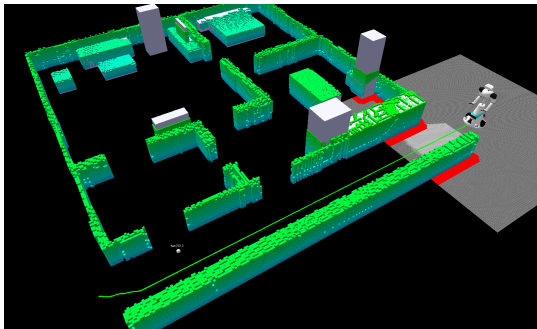
- ▶ Use Octomap for 3D navigation
- ▶ Project columns to 2D costmap and inflate costs and uncertainty for navigation
- ▶ Certainty decays over time instead of known/unknown
 - A wall never moves
 - People are likely to move



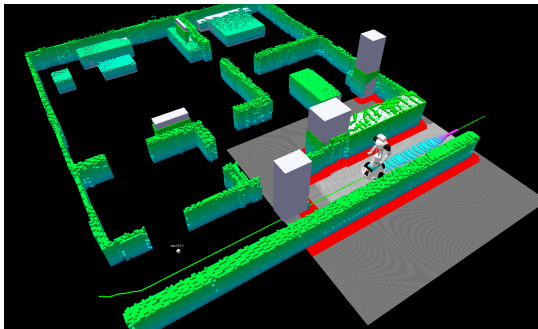
- ▶ Global planner
 - A* Planner
- ▶ Local planner
 - Line collision check
 - Velocities based on safety
 - Assumptions on moving obstacles
 - Desired: DWA/MPC



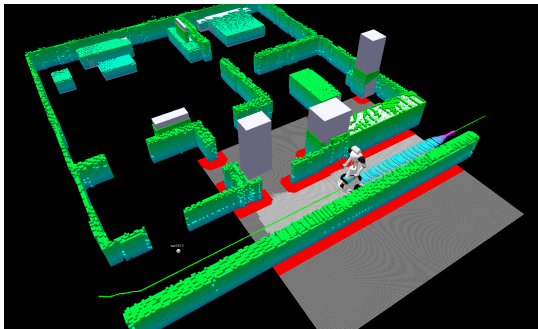
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- ▶ Local planning methods
- ▶ Global planning methods
 - Representations
 - Searching
- ▶ Implementations
- ▶ Further reading: “Motion Planning for Mobile Robots - A Guide”

Questions?